Review of topics in concurrency

What is the smallest value of counter, among those listed, after the threads terminate?

int counter = 0; thread t thread u int cnt; int cnt; for (int i = 0; i < 5; i++) {</pre> for (int i = 0; i < 5; i++) {</pre> 1 5 cnt = counter; cnt = counter; 6 2 counter = cnt + 1;counter = cnt + 1; 3 7 } } 8 4 1. 1

3. 6

4. 10

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The final value of counter is 5 when both threads read counter == 0, one thread proceeds and increments it to 5, and the other thread overwrites the same values up to 5.

But there are schedules where there is an even more destructive interference between the two threads, so that the final value of counter can be as low as 2!

# t'S LOCAL		u'S LOCAL SHARED
$1 \text{ pc}_t : 2 \text{ cnt}_t :$		$pc_u: 6 cnt_u: \perp counter: 0$
2 pc_{t} : 2 cnt_{t} :		$pc_u: 7 cnt_u: 0$ counter: 0
3 pc _t : 3 cnt _t : 0	0	$pc_u: 7 cnt_u: 0$ counter: 0
4 pc _t : 3 cnt _t : 0	0	$pc_u: 6 cnt_u: 0$ counter: 1
5 pc _t : 3 cnt _t : 0	0	$pc_u: 7 cnt_u: 1$ counter: 1
6 pc _t : 3 cnt _t : 0	0	$pc_u: 6 cnt_u: 1$ counter: 2
7 pc _t : 3 cnt _t : 0	0	pc_u : 7 cnt _u : 2 counter: 2
8 pc _t : 3 cnt _t : 0	0	$pc_u: 6 cnt_u: 2$ counter: 3
9 pc _t : 3 cnt _t : 0	0	$pc_u: 7 cnt_u: 3$ counter: 3
10 pc _t : 3 cnt _t : 0	0	$pc_u: 6 cnt_u: 3$ counter: 4
11 pct: 2 cntt: 0	0	$pc_u: 6 cnt_u: 3$ counter: 1
12 pc _t : 2 cnt _t : 0	0	$pc_u: 7 cnt_u: 1$ counter: 1
$13 \text{ pc}_{t} : 3 \text{ cnt}_{t} :$	1	$pc_u: 7 cnt_u: 1$ counter: 1
$14 \text{ pc}_{t} : 2 \text{ cnt}_{t} :$	1	$pc_u: 7 cnt_u: 1$ counter: 2
15 pc _t : 3 cnt _t : 2	2	$pc_u: 7 cnt_u: 1$ counter: 2
16 pc _t : 2 cnt _t : 2	2	$pc_u: 7 cnt_u: 1$ counter: 3
17 pc _t : 3 cnt _t : 3	3	$pc_u: 7 cnt_u: 1$ counter: 3
18 pc _t : 2 cnt _t : 3	3	$pc_u: 7 cnt_u: 1$ counter: 4
19 pc _t : 3 cnt _t : 4	4	$pc_u: 7 cnt_u: 1$ counter: 4
20 done		$pc_u: 7 cnt_u: 1$ counter: 5
21 done		done counter: 2

int n = 0; Semaphore s = new Semaphore(1); // capacity 1 thread t_k int x; s.down(); 1 2 X = N; 3 n = x + 1;s.up(); 4 1. Between 1 and 8

- 2. Between 4 and 8
- 3. Always 4
- 4. Always 8

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int n = 0; Semaphore s = new Semaphore(2); // capacity 2 thread t_k int x; s.down(); 1 2 X = N; 3 n = x + 1;s.up(); 4 1. Between 1 and 8

- 2. Between 4 and 8
- 3. Always 4
- 4. Always 8

int n = 0; Semaphore s = new Semaphore(2); // capacity 2

thread	t _k
--------	----------------

int x;

- 1 S.down();
- 2 X = n;
- 3 n = x + 1;
- 4 s.up();
 - 1. Between 1 and 8
 - 2. Between 4 and 8
 - 3. Always 4
 - 4. Always 8

The value 1 occurs if one thread *t* reads 0 initially, and then waits inside its critical section, while the other threads go through their critical section in mutual exclusion. Then, *t* finishes by writing 1, thus overwriting the increments of all other threads.

```
monitor class CountPrint {
    private Condition isX = new Condition();
    private Condition isY = new Condition();
    public void x()
    { isX.wait(); System.out.print("X"); isY.signal(); }
    public void y()
    { isY.wait(); System.out.print("Y"); isX.signal(); }
}
```

1. A sequence of alternating x and y.

- 2. The first answer, if the monitor uses "signal and wait".
- 3. The first answer, if the monitor uses "signal and continue".
- 4. The program deadlocks.

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- 1. A sequence with at least one x between every pair of Y.
- 2. The first answer, if the monitor uses "signal and wait".
- 3. The first answer, if the monitor uses "signal and continue".
- 4. The program deadlocks.

```
monitor class CountPrint {
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Under "signal and continue", it is possible that two unblocked calls to y() get in the entry queue and then execute one after another.